1. Intruductory rencise

Assume a siriple calculation
$\mathrm{X}=\mathrm{A} \times \mathrm{B} \div \mathrm{C}$
Where $A=10$
$B=20$
$C=30$
To solve this in keyboard mode:-
a) Make sure $S X$ is in $O P E$ mode and that 'Printer Off' $\&$ 'Program Select' Buttons are up.
b) Key $10 \times 20+30=$
C) $S X$ will print $\log$ as shown below

$$
\begin{array}{rr}
\operatorname{Ea} & + \\
\operatorname{Eat} & +
\end{array}
$$

Fig. 1

Now let's try programming this:-
The mode in winch the $S X$ stores programs is 'IEARN' mode
i) To invoke 'LEARN' mode press the 'LRN' key
ii) Depress C, C-ALL to clear the SX's memory

In 'LEARN' mode the $S X$ will automatically print the step number alongside the instruction given.

- You start off the SPnn where $n$ is a character in the range $0-9$ or $A-F$
- 'SP' is short for 'start program' - it denotes the starting point of a program and inserts a special flag so that the start of any program may be easily found.
- SPnn is a double or two-step instruction - after keying the 'SP' a light marked 'UNFIN' will come on to denote that the instruction is unfinished.
- The 'nn' is added to differentiate one program's start point from another's. We will call this one Program 00 (zero zero), so we key Spoo
- Then key the keyboard instructions keyed before ~ this time they will not be executed but will be stored and printed.
- Ihen depress the 'PRJNT' liey to incorporate a print instruction (appears as a diamond)
- Finally key mitio (short foi end of procram ob)

Now depress 'OPE' to put the SX out of 'LEARN' mode and into 'OPEPATE' mode and depress ' $C$ ' to reset the machine.
'Ihen depress 'START' to start the program.
By now, you should have a printout like that shown in Figure 2



```
Fig. 2
```

If something is wrong, you can very easily put the SX back into 'LEARN' mode (but don't depress $C$ and $C$ ALL this time or you will 'wipe' your program) : position yourself to the right step number as shown on the listing by depressing 'STEP SET' and then keying the 3-digit step number of the step to be rekeyed, then rekey the steps as necessary. (Note: the 'Step back' key takes you back to the previous step)

Notice that the answer produced by the 'PRINT' (Diamond) instruction is in 20 column floating-point format. It always appears this way regardless of the setting of the decimal-point wheel. For this reason, it is only used for simple output.

To produce neater, more intelligible output the following technique is uscd:-

First Step - Convert the number from the internal floating-point format to fixed point representation by the 'FIX' instruction:e.g. FIXD FIX5 FIX9

02 Ø2 02
Round Down Round off. Round Up

- to two decimal places

Second Step - Specify a colurn print instruction: COL nn

Where $n$ is the number of columns to be occupied by the printout including sign and decimal point. Naturally the number of columns specified must be large enough to contain the maximum size of number printed with the number of decimal places shown in the 'FIX' Instruction or an error will result.

So to clean up the print out in this example, put the SX into 'LEARN' mode, Step Set 011 and key the following steps:-

FIX5 (keyed as 'FIXnn', '5') (Round Off)
02 (to 2 places)
COL (Column Print)
08 (to 8 places)
EP
00
Then revert to 'OPE' mode, depress ' C ' and then 'START'.
The printout by now should look like that in Figure 3

G11 FING
G1E EE
ESE COL
git Ee
EU15 EF
कtE ge

By now, one monstrows deficioney of the progran should be apparent it will only woric for: vaiues of $\hat{A}$, B \& $C$ of $10,20 \& 30$ : any other values of $A, f$ \& $C$ cquize altaiation to the program.

So what should be done is to allow the operator to key in new values of $A, B \& C$ each time the profanam is run.
How is this to be achieved?
2. Data Entry

Tne SX will stop, turr the 'LNT' light on (and wait for the, operator to key data and press 'SiART' before resuming) under 2 circumstances:-
a) When it encounters an 'SP' instruction in the program (NB this does not happen when a program is called as a subroutine - explained below)
b) When it encounters an 'ENTRY' instruction - obtained by depressing the 'ENT' key and appearing on the listing as 'E'.

| G日g | Et |
| :---: | :---: |
| ¢91 | 5 |
| Q95 | \% |
| 96E | $E$ |
| Egi | $+$ |
| Esc | $E$ |
| פet | $=$ |
| 96 | Fİ |
| Q0\% | E2 |
| QEE | ELL |
| Q19 | 0 |
| -12 | $E$ |
| Q1E | E |

EG.EE
Fig. 4

In Figure 4, the 'SPOO' at the start allows the operator to enter the first number - A. The first ENT allows $B$ to be entered, the second ENT allows $C$ to be entered. The $'=\prime^{\prime}$ calls out the result from the SX's own working registers into the $A$ register for round-off and print out (every algebraic expression has to terminate in an '=' to recall the result).

If the 'Printer-off' button is left up, the $S X$ will $\log$ items entered on the printer; if it is depressed, only explicit print instructions (PRINT, COLUMN - PRINT) will result in print output.

You will notice that after keying the 3 data items, and after printout of the result, the 'ENT' light will again come on. This is because the 'EPOO' statement causes the $S X$ to look for the start of program 00 , and await input again. Thus there is no reason why a program cannot have several EP's, as the EP merely says 'go back to start! A program can obviously have only one 'SP'
3. Introduction to the User of He vies

What if we wanted to store $A, B$ or $C$ for later use in the program? The answer is simple - the instruction 'SMn' (nn is a 2 digit memory no. betwecn 00 and 99) says: 'store the number currently in the A - Register into the nominated memory, leaving the A - register untouched'. When we need to recall the contents of the memory for calculation or printout, we issue the instruction 'RMnn' (short for 'Recall Memory') to bring it back from the nominated memory into the A-Register.
'SMnn' and 'RMnn' are, like the 'SPnn' and 'EPnn' instructions, Double instructions. When, for example, keying in the 'SMnn' instructions in 'LEARJ' Mode, you depress the 'SMnn' key; the step number and 'SM' will be printed on the listing, and the 'UNFINISHED' light will be displayed - the 2 digit memory number is then keyed. Try the example shown in Figure 5.

| bebs Cotet | $9 \mathrm{~g}$ |
| :---: | :---: |
| nger | 51 |
| 6Es | 1 |
| bet | E |
| कge | ? |
| Eete | EP |
| पूर | E |
| 0ege | + |
| -06 | \% |
| Eete | E1 |
| ESII | \% |
| 6 g 5 | Ftr |
| Q61E | PE |
| 0 S 4 |  |
| De15 | $F \mathrm{~F}$ |
| Este | ga |
| פ¢7 | -L |
| bete | Et |
| ghts | LF |
| Deg | LF |
| bect | EF |
| -eme | E |

ig. 5

Here it has been decided to retain $A \& B$ in memories $1 \& 2$. For illustration purposes, the calculation is performed using Memories $1 \& 2$. So, in Figure 5, the nuber (A) entered at the 'SP' is stored in Memory

01 (SMO1), the next (B) is stored in 42 , and the next (C) is left. sitting in the A-Refister as ro do not (in this case) wish to retain it after the calculation wich follovs. Fill and Rollot recall the contents of memories 1 and 2 as the calculation proceeds.

It is very easy to forget the purposes for which you have assigned memories, so it is advisable to make a list of them as you go. This will also mako it very much easier for anyone else reading your program to understand it.

Some other useful memory instructions are:-
CMnn Clear memory nn (for clearing totals)
Emn Add the contents of the A-Register to Nemory nn, leaving the contents of the A register unchanged.

Notice all data movement and calculation has to take place via the A-Register. To copy the contents of Memory 1 to Memory 2 , for instance, the following would be required:-

RM
$\emptyset 1$
SM
$\oint 2$
Also note that the previous contents of the A-Register would be replaced with a copy of Memory 1.

Also note (in Fig. 5 Steps 19-20) the use of the 'LF' (line feed) instruction. This causes the printer to space a line, and is obtained by depressing ' $I / O^{\prime}$ and then ' 0 ' (zero).

In the examples in Figures 4 and 5, it would be very easy for the operator to lose track of when to enter $A$, when to enter $B$, and when to enter $C$. One ideal way of overcoming this is to print instructions on the printer.
How is this to be accomplished?
4. Incorporating punter Messages into the Program

Apart from 'pronpting' the operator, printer messages are also very useful for making the results easier to understand.

To cause a message to be printed in 'Keyboard' or 'Imediate' mode, proceed as follows:-

- Put the SX in 'OPE' mode
for SX100:-
- Depress 'Character Print'
- then depress 'INTnn' followed by the 2 digit code corresponding to the letter required as per the table below (for digits A to $F$, you will see the letters marked under, the 2 leftmost columns of keys, e.g. 'ARC' = A, 'SIN' = B etc)

Repeat for successive characters of the message

- Depress 'Character Print' to terminate the message.

| SX100 Character Table |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Letter | Code | Letter | Code | Letter | Code |
| A | 41 | J | 4A | S | 53 |
| B | 42 | K | 4 B | T | 54 |
| C | 43 | L | 4C | U | 55 |
| D | 44 | M | 4D | V | 56 |
| E | 45 | N | 4 E | W | 57 |
| F | 46 | 0 | 4 F | X | 58 |
| G | 47 | P | 50 | Y | 59 |
| H | 48 | Q | 51 | Z | 5A |
| I | 49 | R | 52 | Space | 20 |

For SX300:-

- Depress 'Character Print'
- Type the message, using the keyboard. Note that the alternative, alphabetic, values of the function keys are displayed below them on the case: e.g. 'ARC' gives the letter 'A' etc.
- Depress 'Character Print' to terminate the message.

To do the same thing under program control, insert the message (in 'LEARN' mode) at the appropriate spot in the program, preceded by 'Character Print', and followed by another 'Character Print' to terminate the message.
.
Note that 'Character Print' appears on the listing as 'CHA'.
Note that the full character set is shown in the SX100 and SX300 manuals. You will see from Figure 6 that our previous example has now had messages inserted so that the operator is told what to enter next
（ENTER A？，ENTER D？，etc）；the answer is preceded by the word＇ANS＇．

| ¢¢ | Sti Lf |
| :---: | :---: |
| Qg | －uF－r |
| Sta | Qte |
| 9EE | －44 |
| ERG | Q4\％ |
| Qer b | Git E |
| GEE GE | －बत7 |
| 9G\％ | णो\％ |
| Qer | \％ |
| E69 | －6 |
| Gib E | 951 \％ |
| EGIL F | MEE E |
| GiE | GES FIGS |
| 6is | G6E |
| ETi4 | GEE E！ |
| QEIE－ | gere |
| GEIE E | Q5： |
| E日可 | GE：+ |
| GEI： | Gcter |
| QTEP号 | EEG |
| QEEGE | GEC |
| GE1 GL | GEE EH |
| Gre | SEE EE |
| QEE LF | SEL $=$ |
| GE4 bth | SEEC $\rightarrow$ |
| P9E5E | QEEE |
| FGEE | SEF GF |
| GET T | EEEG |
| GE | ERE |
| E2G F | Etre |
| 里三甴 | ロ゙ア゙】 |
| ESE | ज6－ |
|  | EqJ Fip |
|  | E日T4 ER |
| EGUE | GTE |
| GE5 | Gty Et |
| ESE EL | G\％LF |
| EGG FTS | Gra |
| 60\％ | ¢9 |
| た6゙G TL | EEFEF |
| 564 5 | Q日 |

Fig． 6

The printout from running the program is shown in Figure 7.

Fit B EOA E

EMTEF Fí IE Q
ETTEF E\％ESGE
EHTEF EG ES GE

Another thing that has had to be done is to avoid the program stopping at Spoo, after the prozram has processed the output for the first set of values. If it did stop at SPOO, after encountering EROO after the first run through the prograw: (EPOO having 'triggered' a search for Spoo), then the 'ENT' iight would cone on without any preceding printout as to what the operator was supposed to enter.

To avoid this, we have to stop the program reaching EPOO, and instead of relying on EFOO to take us back to the start of the program, we instruct the SX (Figure 6, steps 78-79): GTO1 - Meaning 'go to Flag 01' obtained by depressing the 'GO TO nn' key, then keying'01'. We define Flag 01, as being the point before, the printout 'ENTER A?', by keying 'FLAGnn' and 01 at Lines 2-3 (Fig. 6).

Under these circumstances the SPOO and EPOO, at the beginning and end of the program respectively, are no longer required, as the GTOl and FLGØ1 have effectively taken over their function in this case. They are only left in for clarity's sake.
5. Jumins

Quitc often, as in the previous example, the program nceds to JUMP to a point other than one which would be reached naturally. Thesc unconditional jumps are achicved by inserting 60 TO mn at the point in the progra:: at which the jurap is to take place. The 'nn' is a 2 digit code used to definewhere you want to JUiP to, in combination with a FLAG nn instruction. For consistent results, the 'nn' code associated with the FLAGnn must be unique within the program. The 'GO 'O' trizerers a search of memory for a FLAG with the same code. 'GO TO' is obtained by depressing the 'GO TO nn' key, followed by the 2 digits of the flag; it appears on the listing as 'GT'. 'FLAG' is obtained by depressing the 'FLAG nn' key, again followed by the 2 digit code. 'flag' appears on the listing as "FLG'.

Each digit in the code must be in the range $\varnothing-9$, A-F.
So a 'GT55' instruction will trigger a search of memory for a FLG55, if the SX cannot find it, the search will continue over and over again in an endless loop (key ' $C$ ' to terminate).

Using this method of jumping, the SX will be searching for a symbol (FLG + Code) - hence this is known as a symbolic jump.

There is another method - an absolute jump. This is implemented by 1 oading the step number (to be jumped to) into the A-Register, by computation, recalling memories, etc. and then issuing the IOF instruction (Keyed as 'I/O', $\overline{\mathrm{F}}$ ). This is a much faster instruction because the $S X$ does not have to search - it 'knows' which step to go to. However it is inadvisable to use this approach during program development, as any change to a 'destination' step no. through insertion/deletion of other steps would necessitate changing all the IOF sequences pointing to that step. Consequently, it is usual to insert the IOF's only when the program is virtually in final form.
6. Subroutines

Notice from Figure 6 that certain blocks of steps are repeated several times, viz:-

CHA
E
N
$T$ Lines 6-11, 24-29, 42-47
E
R

FIX5
$\emptyset 2$
COL Lines 19-23, 37-41, 53, 56, 73,77
ø8
LF
It is tedious to have to enter these steps over and over again, and also consumptive of steps.

To overcome this problem, we can make use of a technique known as subroutining.

Using this technique, we code a frequently - used routine as a program on its own, beginning with 'SPnn' and ending with 'EPnn'.

When we want to use this routine, we use the instruction 'GO TO SPnn' (short for 'Go to Subprogram') followed by the 2 digit code assigned to the subprogram. The 'GO To SPnn' key prints as GS.

When a subroutine is called via the 'GS' instruction (eg GS21) the SX searches for an 'SP' with the same code - in this case SP21.

When it finds it, it transfers program control to the new subprogram without stopping for input.

When the subprograms 'EP' is encountered (End of Program), the SX, knowing that the subroutine was called as a subroutine (subprogram), transfers program control back to the step after the 'GS' which called the subroutine. Where there are many calls in the same main program to the same subroutine (via several 'GS's' in various parts of the program), the $S X$ keeps track of which 'GS' called the subroutine, so as to be able to return to the correct section of the program, eg:-


Control always transfers back to the step after the 'GS' which called it. Note the difference between the effect of 'EP' in a program called as a 'main' program and 'EP in a program called as a subprogram:-

Main Program: 'EP' causes search for 'SP' Subprogram: 'EP' causes return to calling program The same routine can be used as a main progran and as a subprogram the difference lies in whether it was called via 'GS' or not.

The subroutine can be located anywhere in memory.

| GEE E | $\operatorname{geg} \mathrm{g}$ |
| :---: | :---: |
| tei | 649 |
| G日E PL | cti $\quad 17$ |
| GEE 1 | QRLE QE |
| Ogt Lf | Qers |
| get | 944 |
| bebe Ce | 945 |
| bet | QuE Orth |
| gee itn | W47 |
| ब曰9 | Ete 1 |
| 6te＇ | 649 |
| Qeti CHE | Egto |
| Qute E | WEI GHF |
| の日i三 | dece ec |
| buti | gese |
| betc 6 c | GSEL ET |
| bie | कgse |
| teli 6 | GEE EF |
| Qup E | EST 9 |
| कुt Ef | Qte EF |
| कहEE | Esege |
| QeEI | EEE CH： |
| gere bfi | QES 5 |
| 日SEE E | GES |
| Eet m | QEE T |
| Gese E | Qet E |
| Eete Es | Descir |
| Eet ge | cees |
| cese 5 | QEer Cht |
| Ebeg de | GEGEP |
| 6ese CH | GEEED |
| Qtei ： | Equ |
| 日包 | ET1 |
| QuEs Chat | Gre FIK |
| ESt | StE 5 |
| 605 63 | G日G ELL |
| deve de | 日里明 |
| bus + | EQTE LF |
|  | gra ep |
|  | Ques |

Fi．g． 9

Figure 9 shows the Program in Figure 6 converted to use two subroutines， one of which（SPø2）prints＇ENTER＇the other（SPØ3）does the rounding off and printing．The＇GS＇statements are underlined．The number of steps saved in this case is trivial，but this technique can frequently save，a great many steps and greatly simplify programming．

## 7. Tables \& Dissections

Let us say that we vant to do a sales analysis whereby we key in the territory nuber ( 1 - 50) and then the invoice amount, for all invoices, in such a maner that 50 territory totals are accunulated, for printout after we have finished the last invoice.

Up to now, when using memories we have always specified the memory number to be used. But with 50 possible memory numbers, this. would be impossibly tedious.

Luckily, the $S X$ provides a very easy way round the problem whereby, instead of storing in a memory specified in the program, we can get the program to put the number of the memory to be used into a 'pointer' memory (any memory can be used for this purpose). To specify that the register nominated is to be used as a pointer, we precede the memory reference instruction with 'IND' (indirect) obtained by depressing the 'INDIRECT' key. Contrast these two approaches:-

| A | B |
| :--- | :---: |
|  | (Memory 20 contains 8 ) |
| SM | IND |
| 08 | SM |
|  | 20 |

Both achieva the same effect - in case B the SX sees 'IND' and knows that for the following 'Store Memory' instruction it has to treat the contents of memory 20 as the pointer to where it really has to store the data, instead of storing the data directly into Memory 20.

To do our sales analysis, let's use Memories 1 - 50 for the 50 territory totals, and 51 as storage for the territory number. Then, in order to ensure that the invoice total is accumulated in the correct memory, we just designate 51 as a 'pointer' memory.

GUG
巴gi E
get i

BE FL
BEL EL
ges $\angle F$
DGE V明
बिद $T$
Qgeg E
Gge
ELG
日连 I
EtET
EETE
BEL F
EsIE
REE
EET F
GIGE
QED ：

GEI FTM
ERE
GEE ER
ERL E5
EEDE EHA
EEE
घृ
GEES
BECGF
GE $L$
EDEJ E
gere
EES $\%$
EESL EHF
Gezs E
ESEF FTYC
GEG ED
ESE ERL
QESO 6
ELLE IHL
E41 E 4
GUE E1
GR4 $2 T$
कीLL
EGUS EF
GLE GE

$$
\because 10
$$

Figure 10 shows a program to do this: the operator is asked for the territory number, and then the sales, which are accumulated in the appropriate territory total in nemories 1-50. Notice the useof the 'Fl' instruction at line 2 to clear all memorics - this is obtained by keying 'mistnn', F, 1.

At line 44, the program goes back to line 3 (after the F1 instruction) for the next entry, in anendless loop.

| Ere | $\pm$ | Gbest | 1.9 |
| :---: | :---: | :---: | :---: |
| TEFEITGPE | E | Brtes | .6 |
| TEFELTME? | 3 | GELES | E.ED |
| TEFETAEPG | 4 | geter | 4.65 |
| TEFEITCRE | - | GREE? | 1.6 |
| TEFEITCRE | $E$ | GRLES | .en |
| TEREITARM | $\equiv$ | GHLES | E.EE |
| TEARTTEEV | 4 | GFLEEC | 4.8 |
| TEFAITCRG |  |  |  |


| E | Fets |
| :---: | :---: |
| 4 | froe |
| $E$ | Fte |
| 0 | Pbel |
| $\underline{\square}$ | Ftes |

Figure 11 shows the printout of the program when run('OPE' mode, ' $\mathrm{C}^{\prime}$, \& 'START' - 'PRINTER OFY' down).

After the last entry, it was necessary to depress ' C ' to stop the program, and put the 'PRINTER OFF' button up, in order to get the printout shown by manally recalling memories 1 - 5 (tedious for 50!) To recall a memory, you simply key 'RMnn' followed by the memory no:Ø1, ø2, Ø3 etc.

Wouldn't. it be nice to have this printout occur automatically? To achieve this, you have to master 2 new skills, - Use of Conditions, and Loops
8. Use of Conditions

It is often necessary for a program to make a logical decision. In the case discussed above, it wouid be nice to have the machine detect whether the last entry has been keyed.

This might be achieved by use of the 'IF FNT' test. This test checks whether the operator has entcred anything in response to the 'ENT' command. It is entered by keying:-

IF GO TO nn
ENTRY
dd
('dd' is the two-digit code of the flag to which the program is to go if something (even a zero) has been keyed)

If appears on the listing as, eg:-
IFE
$\emptyset 5$
So to terminate the entries, the operator would depress 'START' without keying anything.

Other conditional Tests are available for testing the contents of the A-Register after an arithmetic operation:-

Purpose
Is A Non-Zero?
Is A Positive or Zero
Is A Negative

Keyed As
If GO TO, $=$, dd
If GO TO, + , dd
If GO TO, -, dd

## Lists As

IFNZ
dd
dF+
Td

So to test whether the number contained in Memory 02 is less than or equal to that in $\emptyset 1$, and if so to go to FLAG 80 , the following could be used:-

RM
$\emptyset 1$
RM
$\emptyset 2$
$=$ (don th forget to recall result with $\quad=^{\prime \prime}$ )
IF+
80
If the number in Memory 2 is greater than that in Memory 1 , (test is not true), control will 'fall through' the 'IF' and continue at the following step.

Other forms of teat are listed in the $S X$ Programmer's Manual. The instructions to test whether the operator keyed anything are shown in Figure 12, Steps 18-20.

```
अघQ 5P
G%G
Gav & 
GGG FLG
CELTE
gem LF
QGE EHB
GEGT
GEGE
GGF
GEE F
Evi I
GEIE T
BEIS !
ELLF
#55
GEF
BET EAH
G1:
```




```
GED ET
GEE GO ro-6o 70. dq
GEF FLE
&54 %
```




```
GgeT FIGG
EESE EE
GEG ELL
GE 5
DEZ1 EHF
EGZO
&6
60%%4
EEES F
BEEL
807 E
EGGE
EOG
GGLE EHF
EG4工 E
GGEFINE
#E4马 5
G4L OL
GBGC
GGLF IT[ GCCUNIGAMTE IN:CICE
```



```
EtS 5D
```

Q4G 1


9. Loops

In the Sales Analysis case under discussion, we want to step through the pointer through the values $1-50$, so that we can print out the corresponding totals.

To do this we can construct a simple loop, Using Memory 51 as a counter:-


The steps to do this (one way) are shown on 1 ines $83-97$ and 126-127 Other refinements added are as follows:-
a) In lines 100-104 we test if the Territory total is Non-Zero if so, well print it, otherwise well skip that territory and go on to the next.
b) In lines 115-116 we make use of another handy print formatting instruction - 'SPACEnn' (prints as sideways arrow) which means 'space nn positions' in this case 4 print positions.

The run of the program is shown in Figure 13.

| OEFETTGEP | $\pm$ | GEET |  |
| :---: | :---: | :---: | :---: |
| TEFEITEA | 2 | C－LES | E．ge |
| TEFEITCET | 3 | GREG？ | E．ET |
| TEFEITGRU | 4 | ELLES | 4.00 |
| TEFEITGEE | 1 | GLEET | 1．09 |
| TEFEITCRE | ： | GREST | E．ts |
| TEFEITEA？ | 3 | GFLEC？ | E．En |
| TEFETTFG\％ | 4 | GFLES？ | 4.60 |
| TEFEITCEA | 2 | GLES | 19．60 |
| TEFEITES | 35 | GFLES＇ | 三E．LE |
| TEFEITCE＇ | 5 | GFLES | 念时 |

TEFEITGFG TGTFL GFLEE


## 10. Function Keys

Frequently the operator necds to be able to call up a special routine such as printing totals (as in the provious example), entering a credit, doing an error corrcction, etc.

A Function Key facility is provided for this purpose:- when the 'PROGKil SELECT' button is down, the following 6 keys change their function and, in OPE mode, become function keys:-

| ARC | (A) | $\mathrm{e}_{2}^{\mathrm{X}}$ | (F) |
| :--- | :--- | :--- | :--- |
| SIN | (B) | $\mathrm{a}^{\mathrm{C}}$ | (U) |
| COS | (C) |  |  |
| TAN | (D) |  |  |
| a | (E) |  |  |

Deprossion of a function key causes the $S X$ to look for a routine starting with 'SP' and then a reserved code according to the following:-
Key Code

| A | 8 A |
| :--- | :--- |
| B | 8 B |
| C | 8 C |
| D | 8 D |
| E | 8 E |
| F | 8 F |
| U | 89 |

So a routine to be activated by the 'C' key will start:-
SP
8C (etc)
For this to occur, the 'Program Select' button must be down, and the SX must either be idle cr in the 'ENT' state (awaiting entry).

A special keyboard overlay is available to remind the operator which function key performs what function.

Another way of calling up a special function when in 'OPE' mode is to depress 'GO TO SPnn' followed by the routine's 2 digit code.

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-23-
$$

10. Conclusion

Hopefully, this has served as a primer to the use of the $S X$; it is intended to be read in conjunction with the SX Programuing Manuals

In particular, the following topics are not covered in this Introduction:-

- Use of Check Node for Inserting and Deleting steps
- Use of debug Mode for Debugging programs
- Scientific functions
- Use of Magnetic Card \& Cartridge
- Splitting Memories
- Full SX Instruction Set

For these consult the following Canon Publications:-
SX Programming manual
SX Programmable Calculator Instructions
SX Scientific Functions Instructions

